

QUANTUM CHANNELS II: DATA-PROCESSING, RECOVERY CHANNELS,

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0. Introduction: Quantum Channels

→ $\mathcal{H}_A, \mathcal{H}_B$: finite-dim. Hilbert spaces

$\mathcal{B}(\mathcal{H}_A), \mathcal{B}(\mathcal{H}_B)$: algebras of linear op's acting on $\mathcal{H}_A, \mathcal{H}_B, \dots$

→ A quantum state $\rho_A \in \mathcal{B}(\mathcal{H}_A)$ is a positive semi-definite op of unit trace:

$$\rho_A \geq 0, \quad \text{tr} \rho_A = 1$$

→ If $\text{rk} \rho_A = 1$, then ρ_A is called a pure state, then exists $|\rho_A\rangle \in \mathcal{H}_A$,

$$\langle \rho_A | \rho_A \rangle = 1, \quad \text{s.t.} \quad \rho_A = |\rho_A\rangle\langle \rho_A|.$$

→ A quantum channel is a linear map $\mathcal{N}: \mathcal{B}(\mathcal{H}_A) \rightarrow \mathcal{B}(\mathcal{H}_B)$ that is

→ completely positive: $\mathcal{N} \otimes \text{id}_n$ is positive $\forall n \in \mathbb{N}$

$$(CP) \quad (\text{T positive: } (X \geq 0 \Rightarrow \mathcal{T}(X) \geq 0),$$

$$\text{id}_n: \mathcal{B}(\mathbb{C}^n) \rightarrow \mathcal{B}(\mathbb{C}^n), X \mapsto X)$$

→ trace-preserving: $\text{tr} \mathcal{N}(X_A) = \text{tr} X_A \quad \forall X_A \in \mathcal{B}(\mathcal{H}_A)$

$$(TP)$$

→ $\mathcal{N}: \mathcal{B}(\mathcal{H}_A) \rightarrow \mathcal{B}(\mathcal{H}_B)$ is CPTP iff

$$(\mathcal{N}: A \rightarrow B)$$

↑ $|B| \times |A|$ matrices

i) (Kraus representation): $\exists K_i: \mathcal{H}_A \rightarrow \mathcal{H}_B$ s.t.

$$\mathcal{N}(X_A) = \sum_i K_i X_A K_i^\dagger \quad (\text{CP})$$

$$\text{and } \sum_i K_i^\dagger K_i = \mathbb{1}_A.$$

ii) (isometric picture): \exists isometry $V: \mathcal{H}_A \rightarrow \mathcal{H}_B \otimes \mathcal{H}_E$

for some auxiliary Hilbert space \mathcal{H}_E ("environment") s.t.

$$\mathcal{N}(X_A) = \text{tr}_E V X_A V^\dagger$$

(V is an isometry $\Leftrightarrow V^\dagger V = \mathbb{1}_A$),

$$\text{tr}_E \gamma_{BE} = \sum_{i=1}^{|\mathcal{E}|} (\mathbb{1}_B \otimes \langle i |_E) \gamma_{BE} (\mathbb{1}_B \otimes |i\rangle_E)$$

where $\{|i\rangle\}_{i=1}^{|\mathcal{E}|}$ is an ONS for \mathcal{H}_E ;

$$|\mathcal{E}| = \dim \mathcal{H}_E. \quad \left. \begin{array}{l} \\ \downarrow \mathcal{H}_E \cong \mathcal{H}_B \end{array} \right\}$$

iii) (unitary picture): \exists unitary $U: \mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_{E'} \rightarrow \mathcal{H}_{ABE'}$

and a state $|\varphi\rangle \in \mathcal{H}_B \otimes \mathcal{H}_{E'}$ s.t.

$$\mathcal{N}(X_A) = \text{tr}_{AE'} U (X_A \otimes |\varphi\rangle\langle\varphi|_{BE'}) U^\dagger$$

.) The Choi operator τ_{AB}^N of a linear map $N: A \rightarrow B$ is defined as

$$\tau_{AB}^N = (\text{id} \otimes N)(\gamma_{AA'})$$

where $|\gamma\rangle_{AA'} = \sum_{i=1}^{|A|} |i\rangle_A \otimes |i\rangle_{A'}$ ($\{|i\rangle\}_{i=1}^{|A|}$ is a basis for $\mathcal{X}_A \cong \mathcal{X}_{A'}$)

is an unnormalized maximally entangled state.

$N: A \rightarrow B$ is .) CP iff $\tau_{AB}^N \geq 0$;

.) TP iff $\text{tr}_B \tau_{AB}^N = \mathbb{1}_A$.

τ_{AB}^N uniquely defines $N: A \rightarrow B$ via the Choi-Jamiołkowski

isomorphism: $N(X_A) = \text{tr}_A \left[\tau_{AB}^N (X_A^T \otimes \mathbb{1}_B) \right] \quad \forall X_A \in \mathcal{B}(\mathcal{X}_A)$

.) Let $N: A \rightarrow B$ be a linear map. The adjoint map $N^\dagger: B \rightarrow A$

is defined via $\langle N^\dagger(\eta_B), X_A \rangle = \langle \eta_B, N(X_A) \rangle$

$\langle X, \eta \rangle := \text{tr}(X^\dagger \eta) \quad \forall X_A \in \mathcal{B}(\mathcal{X}_A), \eta_B \in \mathcal{B}(\mathcal{X}_B)$

Hilbert-Schmidt
inner product

.) A linear map $N: A \rightarrow B$ is .) CP iff N^\dagger is CP

.) TP iff N^\dagger is unital:

$$N^\dagger(\mathbb{1}_B) = \mathbb{1}_A.$$

.) For a CP map $\mathcal{N} = \sum_i K_i \cdot K_i^\dagger$, the adjoint map is

$$\mathcal{N}^\dagger = \sum_i K_i^\dagger \cdot K_i$$

Example: $\mathcal{N} = \text{tr}_B: A \otimes B \rightarrow A$

$\{|i\rangle\}_{i \in \mathcal{I}}$ ONB for \mathcal{H}_B

$$\text{tr}_B: X_{AB} \mapsto \sum_i \underbrace{(\mathbb{1}_A \otimes \langle i|_B)}_{= K_i} X_{AB} (\mathbb{1}_A \otimes |i\rangle_B)$$

$$\begin{aligned} \mathcal{N}^\dagger(\gamma_A) &= \sum_i K_i^\dagger \gamma_A K_i = \sum_i (\mathbb{1}_A \otimes |i\rangle_B) \gamma_A (\mathbb{1}_A \otimes \langle i|_B) \\ &= \underline{\gamma_A \otimes \mathbb{1}_B}. \end{aligned}$$

1. Relative entropy and data-processing

§ 1.1 Motivation: quantum state discrimination

Assume you're given one of two quantum states ρ_0, ρ_1 with equal probability $\frac{1}{2}$.

Goal: decide which one you get

Strategy: perform a measurement on the unknown state giving you a guess for the correct answer.

For us: measurement is described by a positive operator-valued

measure (POVM): $\{\Lambda_i\}_{i=1}^k$, $\Lambda_i \geq 0$, $\sum_i \Lambda_i = \mathbb{1}_X$

for a given state ρ , POVM gives the outcome "i"

with probability $\text{tr}(\Lambda_i \rho) = p_i$.

Here, we use a POVM $\Lambda = \{\Lambda_0, \Lambda_1\}$ ($\Lambda_1 = \mathbb{1} - \Lambda_0$, $\Lambda_0 \geq 0$)

outcome "0" corresponds to ρ_0 w.p. $\text{tr}(\Lambda_0 \rho)$

"1" — " — ρ_1 w.p. $\text{tr}(\Lambda_1 \rho)$ $\rho \in \{\rho_0, \rho_1\}$

Success probability of correctly identifying the state:

$$P_{\text{succ}}(\Lambda) = \frac{1}{2} \underbrace{P_r(\rho_0 | \rho_0)}_{\text{tr} \Lambda_0 \rho_0} + \frac{1}{2} \underbrace{P_r(\rho_1 | \rho_1)}_{\text{tr} \Lambda_1 \rho_1}$$

$$\max_{\Lambda = \{\Lambda_0, \Lambda_1\}} \leftrightarrow \max_{0 \leq \Lambda_0 \leq \mathbb{1}}$$

$$= \frac{1}{2} (\text{tr} \Lambda_0 \rho_0 + \text{tr} \Lambda_1 \rho_1) \xrightarrow{\Lambda_1 = \mathbb{1} - \Lambda_0} \text{tr} (\mathbb{1} - \Lambda_0) \rho_1$$

$$= \frac{1}{2} (\text{tr} \Lambda_0 \rho_0 + 1 - \text{tr} \Lambda_0 \rho_1) = 1 - \text{tr} \Lambda_0 \rho_1$$

$$= \frac{1}{2} (1 + \text{tr} \Lambda_0 (\rho_0 - \rho_1))$$

Want to max. $P_{\text{succ}}(\Lambda)$ with respect to POVM $\Lambda = \{\Lambda_0, \Lambda_1\}$

$$P_{\text{succ}}^* = \max_{0 \leq \Lambda_0 \leq \mathbb{1}} P_{\text{succ}}(\Lambda) = \frac{1}{2} (1 + \max_{0 \leq \Lambda_0 \leq \mathbb{1}} \text{tr} \Lambda_0 (\rho_0 - \rho_1))$$

$$P_{\text{succ}} = \max_{0 \leq \Lambda \leq \mathbb{1}} P_{\text{succ}}(\Lambda) = \frac{1}{2} \left(1 + \max_{0 \leq \Lambda \leq \mathbb{1}} \text{tr} \Lambda (\rho_0 - \rho_1) \right)$$

→ trace norm $\|X\|_1 = \text{tr} \sqrt{X+X^\dagger} = \sum_i s_i(X)$

↑ singular values of X
 \Leftrightarrow eigenvalues of $|X| = \sqrt{X+X^\dagger}$

→ trace distance of two states ρ_0, ρ_1 : $\frac{1}{2} \|\rho_0 - \rho_1\|_1$

Lemma 1 Let ρ_0, ρ_1 be quantum states, then

$$\frac{1}{2} \|\rho_0 - \rho_1\|_1 = \max_{0 \leq \Lambda \leq \mathbb{1}} \text{tr} \Lambda (\rho_0 - \rho_1)$$

Proof: $\rho_0 - \rho_1$ is Hermitian: $\rho_0 - \rho_1 = \sum_i \lambda_i |i\rangle\langle i|$

where: $\lambda_i \in \mathbb{R}$, $\langle i|j\rangle = \delta_{ij}$.
 (eigenvalues) (eigenvectors)

$$\left. \begin{aligned} \text{Define: } P &= \sum_{i: \lambda_i \geq 0} \lambda_i |i\rangle\langle i| \geq 0 \\ Q &= \sum_{i: \lambda_i < 0} (-\lambda_i) |i\rangle\langle i| \geq 0 \end{aligned} \right\} \rho_0 - \rho_1 = P - Q$$

$$\|\rho_0 - \rho_1\|_1 = \text{tr} |\rho_0 - \rho_1| = \text{tr} |P - Q| = \text{tr} P + \text{tr} Q = 2 \text{tr} P$$

$$|A| = \sqrt{A^\dagger A}$$

$$\text{tr}(P - Q) = \text{tr}(\rho_0 - \rho_1) = \text{tr} \rho_0 - \text{tr} \rho_1 = 0$$

$$\Rightarrow \text{tr} P = \text{tr} Q$$

$$\Rightarrow \frac{1}{2} \|S_0 - S_1\|_1 = \text{tr } P$$

relate this to $\text{tr } \Lambda (S_0 - S_1)$ for arbitrary $0 \leq \Lambda \leq \mathbb{1}$:

$$\text{tr } \Lambda (S_0 - S_1) = \text{tr } \Lambda (P - Q)$$

$$\begin{aligned} Q \geq 0 \\ \Lambda \geq 0 \\ \leq \text{tr } \Lambda P \end{aligned}$$

$$\Lambda \leq \mathbb{1} \leq \text{tr } P = \frac{1}{2} \|S_0 - S_1\|_1$$

$$\Rightarrow \max_{0 \leq \Lambda \leq \mathbb{1}} \text{tr } \Lambda (S_0 - S_1) \leq \frac{1}{2} \|S_0 - S_1\|_1$$

remains to be shown: $\exists \Lambda$ achieving maximum.

Set $\Pi_P = \sum_{i: \lambda_i \geq 0} |\lambda_i\rangle\langle\lambda_i|$, projector onto support of P :

$$\Pi_P^2 = \Pi_P, \quad \Pi_P \geq 0, \quad \Pi_P \leq \mathbb{1}, \quad \Pi_P P \Pi_P = P, \quad \Pi_P Q \Pi_P = 0$$

$$\text{tr } \Pi_P (S_0 - S_1) = \text{tr } \Pi_P (P - Q)$$

$$= \underbrace{\text{tr } \Pi_P P}_{= P} - \underbrace{\text{tr } \Pi_P Q}_{= 0} = \text{tr } P = \frac{1}{2} \|S_0 - S_1\|_1.$$

□